

A multi-stage stochastic programming approach for remanufacturing planning under uncertainty

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PGMO Days 2016

Plan

- 1 Problem presentation
- 2 Literature review
- 3 Stochastic problem: model
- 4 Stochastic problem: resolution
- 5 Conclusion and perspectives

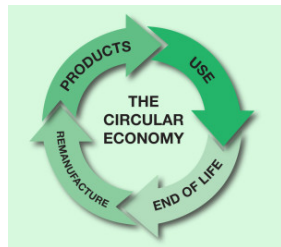
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Remanufacturing

Context

Transition towards a circular economy



Remanufacturing

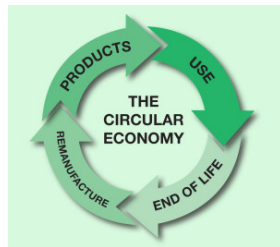
Context

Transition towards a circular economy

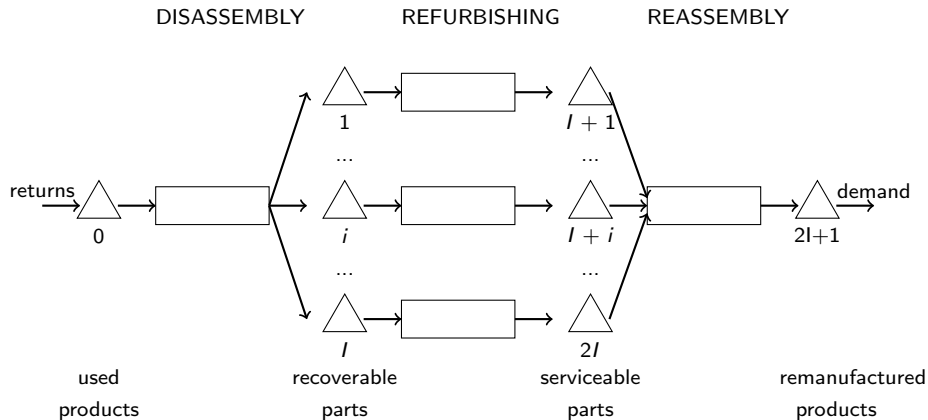
Remanufacturing: a way of rehabilitating used products

Obtain new (or like-new) products

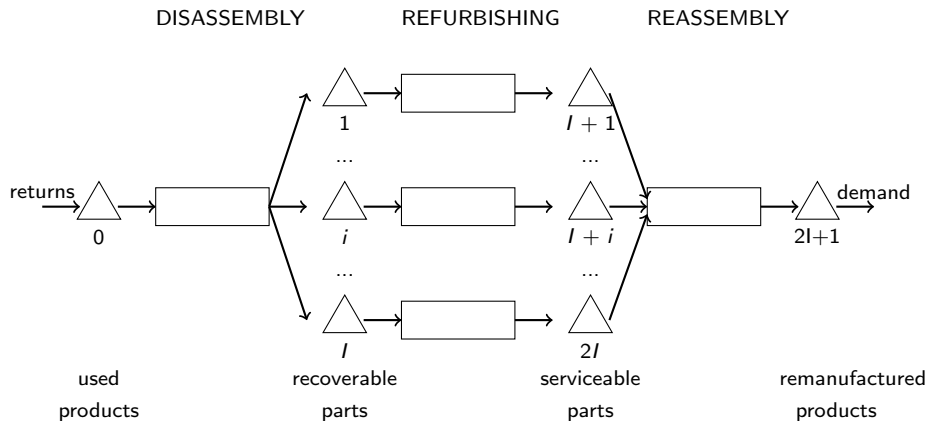
- from end-of-life products returned by consumers
- by replacing or refurbishing damaged components



A three-echelon remanufacturing system



A three-echelon remanufacturing system



Industrial applications: mobile phones, electrical equipment...

[Jayaraman, 2006], [Franke et al, 2006], [Han et al, 2013], [Ahn et al, 2011]

Production planning & lot-sizing

Production planning

Over a multi-period horizon, decide:

- how many used products to disassemble,
- how many parts to refurbish,
- how many remanufactured products to reassemble,

so as to satisfy the demand for remanufactured products.

Production planning & lot-sizing

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Costs

- Fixed production setup costs
⇒ A small number of large-size production lots
- Inventory holding costs
⇒ A large number of small-size production lots

Production planning & lot-sizing

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⇒ A small number of large-size production lots
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Lot-sizing

Plan production so as to optimize the trade-off between setup and inventory holding costs

Assumptions

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- 1 a single type of returned/remanufactured product

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- 1 a single type of returned/remanufactured product
- 2 homogenous quality of the returned products

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- ① a single type of returned/remanufactured product
- ② homogenous quality of the returned products
- ③ deterministic yield of the disassembly and refurbishing processes

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- 1 a single type of returned/remanufactured product
- 2 homogenous quality of the returned products
- 3 deterministic yield of the disassembly and refurbishing processes
- 4 uncapacitated production processes

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- ③ deterministic yield of the disassembly and refurbishing processes
- ④ uncapacitated production processes

Deterministic problem: model

Notation

- planning periods: $t = 1 \dots T$
- products: $i = 0 \dots 2l + 1$
- production processes: $p = 0 \dots l + 1$

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- α_i : quantity of product i needed to produce one remanufactured product
 $\alpha_0 = 1$; $\alpha_i = \alpha_{i+1}$, $\forall i = 1 \dots l$; $\alpha_{2l+1} = 1$

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- r_t : quantity of used products returned in t
- d_t : demand for remanufactured products in t

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Decision variables

- $Y_p^t \in \{0, 1\}$: setup variable for process p in t
- X_p^t : production quantity for process p in t
- I_i^t : inventory level for product i at the end of t
- L^t : lost sales of remanufactured product in t

Deterministic problem: formulation

MILP formulation

$$Z^* = \min \sum_{t=1}^T \left[\sum_{p=0}^{l+1} s_p Y_p^t + \sum_{i=0}^{2l+1} h_i I_i^t + l L^t \right]$$

$$I_0^t = I_0^{t-1} + r^t - X_0^t \quad \forall t$$

$$I_i^t = I_i^{t-1} + \alpha_i X_0^t - X_i^t \quad \forall t, \forall i = 1 \dots l$$

$$I_i^t = I_i^{t-1} + X_{i-1}^t - \alpha_{i-1} X_{i+1}^t \quad \forall t, \forall i = l + 1 \dots 2l$$

$$I_{2l+1}^t = I_{2l+1}^{t-1} + X_{l+1}^t - d^t + L^t \quad \forall t$$

$$X_p^t \leq M Y_p^t \quad \forall t, \forall p = 0 \dots l + 1$$

$$Y_p^t \in \{0, 1\} \quad \forall t, \forall p = 0 \dots l + 1$$

Deterministic problem: formulation

MILP formulation

	Setup	Inventory	Lost sales
$Z^* = \min \sum_{t=1}^T \left[\sum_{p=0}^{I+1} s_p Y_p^t + \sum_{i=0}^{2I+1} h_i I_i^t + l L^t \right]$			
$I_0^t = I_0^{t-1} + r^t - X_0^t$			$\forall t$
$I_i^t = I_i^{t-1} + \alpha_i X_0^t - X_i^t$		Inventory balance	$\forall t, \forall i = 1 \dots I$
$I_i^t = I_i^{t-1} + X_{i-1}^t - \alpha_{i-1} X_{i+1}^t$			$\forall t, \forall i = I + 1 \dots 2I$
$I_{2I+1}^t = I_{2I+1}^{t-1} + X_{I+1}^t - d^t + L^t$		Setup- production	$\forall t$
$X_p^t \leq M Y_p^t$			$\forall t, \forall p = 0 \dots I + 1$
$Y_p^t \in \{0, 1\}$			$\forall t, \forall p = 0 \dots I + 1$

Research objectives

Uncertain input data

Difficulty to forecast the future returns and demand

→ r_t and d_t = stochastic parameters

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Assumption

Known probability distribution of the stochastic parameters

Research objectives

Uncertain input data

Difficulty to forecast the future returns and demand

→ r_t and $d_t =$ stochastic parameters

Assumption

Known probability distribution of the stochastic parameters

Objective

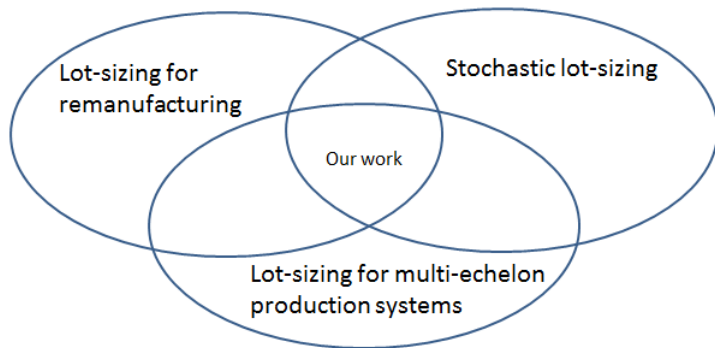
Develop a stochastic programming approach for the problem

- modeling: a multi-stage approach based on a scenario tree
- solving: polyhedral approach, Cut & Branch algorithm

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Literature review



Our contributions

1 Lot-sizing for remanufacturing

Improvement of the problem modeling:

- Multi-echelon system
- Stochastic input data

2 Stochastic lot-sizing

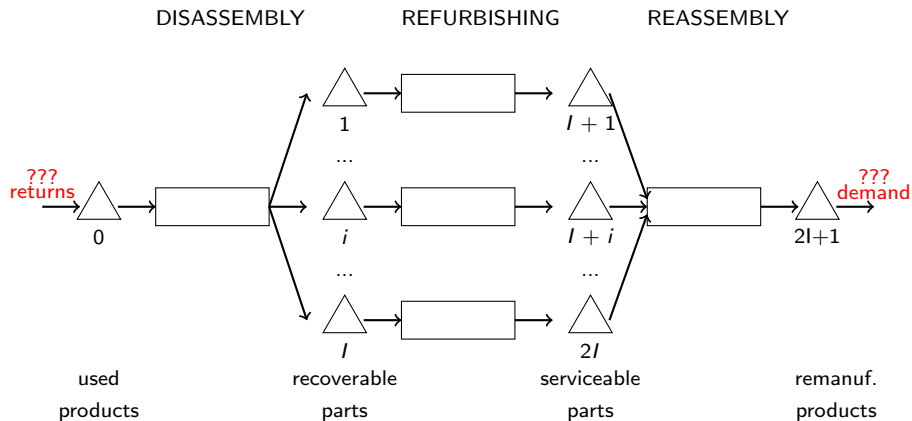
Extension of an existing polyhedral solution approach

- from a single-echelon single-product problem
- to a multi-echelon multi-product problem

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Uncertain returns and demand

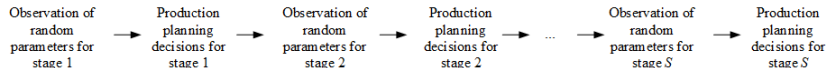


Multi-stage stochastic programming

Assumption

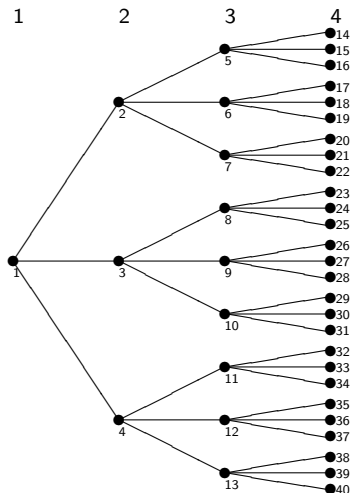
- Not all decisions have to be made before uncertainty realization.
- Some can be postponed to a future point in time.

Dynamic multi-stage decision process



Uncertainty representation

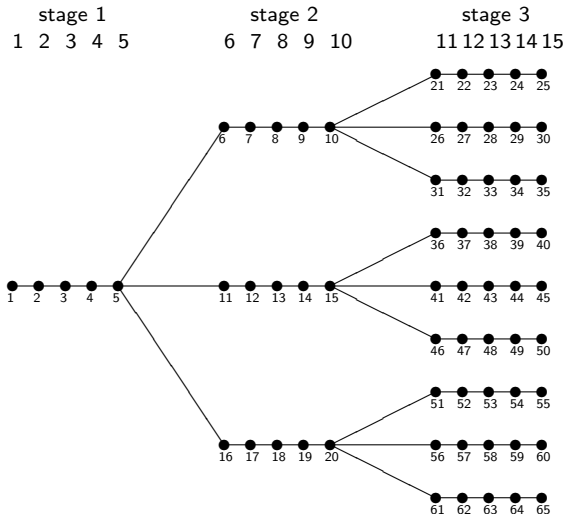
Scenario tree: 1 stage = 1 period



- A node n = a realization of the stochastic returns and demand (r_n, d_n) in the corresponding period
- New information on uncertainty available in each period
→ Every node has several children.

Uncertainty representation

Scenario tree: 1 stage = π periods



- A node n = a realization of the stochastic returns and demand (r_n, d_n) in the corresponding period
- New information on uncertainty available every π periods
 → Only a subset of nodes have several children.

[Zanjani *et al*, 2010]

Problem formulation

Notation

- $n = 1 \dots N$: set of nodes
- π_n : probability of node n
- $a(n)$: predecessor of node n in the scenario tree
- r^n : quantity of used products returned in n
- d^n : demand for remanufactured products in n

Problem formulation

Notation

- $n = 1 \dots N$: set of nodes
- π_n : probability of node n
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Decision variables

- $Y_p^n \in \{0, 1\}$: setup variable for process p in n
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- I_i^n : inventory level for product i at the end of n
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Problem formulation

MILP formulation

Expected setup/inventory/lost sales costs

$$Z^* = \min \sum_{n=1}^N \pi_n \left[\sum_{p=0}^{I+1} s_p Y_p^n + \sum_{i=0}^{2I+1} h_i I_i^n + l L^n \right]$$

$$I_0^n = I_0^{a(n)} + r^n - X_0^n \quad \forall n$$

$$I_i^n = I_i^{a(n)} + \alpha_i X_0^n - X_i^n \quad \text{Inventory balance} \quad \forall n, \forall i = 1 \dots I$$

$$I_i^n = I_i^{a(n)} + X_{i-1}^n - \alpha_{i-1} X_{i+1}^n \quad \forall n, \forall i = I + 1 \dots 2I$$

$$I_{2I+1}^n = I_{2I+1}^{a(n)} + X_{I+1}^n - d^n + L^n \quad \text{Setup-production} \quad \forall n$$

$$X_p^n \leq M Y_p^n \quad \forall n, \forall p = 0 \dots I + 1$$

$$Y_p^n \in \{0, 1\} \quad \forall n, \forall p = 0 \dots I + 1$$

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$$I_0^n = I_0^{a(n)} + r^n - X_0^n \quad \text{Dependent demand} \quad \forall n$$

$$I_i^n = I_i^{a(n)} + \alpha_i X_0^n - X_i^n \quad \forall n, \forall i = 1 \dots l$$

$$I_i^n = I_i^{a(n)} + X_{i-1}^n - \alpha_{i-1} X_{i+1}^n \quad \forall n, \forall i = l + 1 \dots 2l$$

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Solution

Echelon stock = Total inventory of an item in the system, as such or as a component of other items

Problem formulation

MILP 'echelon stock' formulation

$$Z^* = \min \sum_{n=1}^N \pi_n \left[\sum_{p=0}^{l+1} s_p Y_p^n + \sum_{i=0}^{2l+1} eh_i EI_i^n + lL_n \right]$$

$$EI_0^n = EI_0^{a(n)} + r^n - \alpha_0(d^n - L^n) \quad \forall n$$

$$EI_i^n = EI_i^{a(n)} + \alpha_i X_0^n - \alpha_i(d^n - L^n) \quad \forall n, \forall i = 1 \dots l - 1$$

$$EI_{p+l}^n = EI_{p+l}^{a(n)} + \alpha_{p+l} X_p^n - \alpha_{p+l}(d^n - L^n) \quad \forall n, \forall p = 0 \dots l + 1$$

$$X_p^n \leq MY_p^n \quad \forall n, \forall i = 0 \dots l + 1$$

$$EI_0^n - \frac{EI_1^n}{\alpha_1} \geq 0 \quad \forall n$$

$$EI_i^n - EI_{i+1}^n \geq 0 \quad \forall n, \forall i = 1 \dots l$$

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$$X_p^n \leq MY_p^n \quad \forall n, \forall i = 0 \dots l + 1$$

$$EI_0^n - \frac{EI_1^n}{\alpha_1} \geq 0 \quad \begin{array}{l} l + 1 \text{ independent} \\ \text{subproblems} \end{array} \quad \forall n$$

$$EI_i^n - EI_{i+1}^n \geq 0 \quad \forall n, \forall i = 1 \dots l$$

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$$Y_p^n \in \{0, 1\} \quad \forall n, \forall p = 0 \dots l + 1,$$

Subproblems

Subproblem for production process p

$$Z^* = \min \sum_{n=1}^N \pi_n \left[s_p Y_p^n + e h_{p+1} EI_i^n + l L_n \right]$$

$$EI_{p+1}^n = EI_{p+1}^{a(n)} + \alpha_{p+1} X_p^n - \alpha_{p+1} (d^n - L^n) \quad \forall n$$

$$X_p^n \leq M Y_p^n \quad \forall n$$

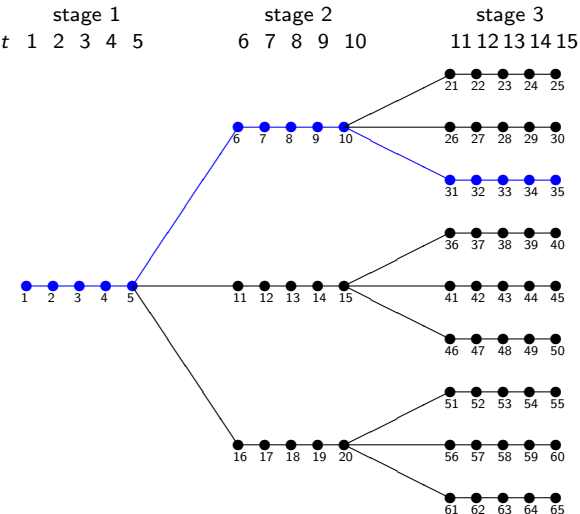
$$Y_p^n \in \{0, 1\} \quad \forall n,$$

→ Single-item single-echelon single-resource lot-sizing with lost sales

→ But on a tree !

Subproblem formulation strengthening

Path inequalities



Subproblem formulation strengthening

Valid inequalities

Let $k \in \{0, \dots, N\}$ be a non-leaf node and $l \in \mathcal{L}(k)$ be a leaf.

Let $U \in \mathcal{P}(k, l)$.

$$EI_{p+l}^k \geq \sum_{n \in U} [d^n (1 - \sum_{\nu \in \mathcal{P}(k, n)} Y_p^\nu - L^n)]$$

Subproblem formulation strengthening

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$$EI_{p+l}^k \geq \sum_{n \in U} [d^n (1 - \sum_{\nu \in \mathcal{P}(k, n)} Y_p^\nu - L^n)]$$

where:

- $\mathcal{L}(n)$: set of leaves in the subtree originated from n
- $\mathcal{P}(n, m)$: path between node n and m in the scenario tree

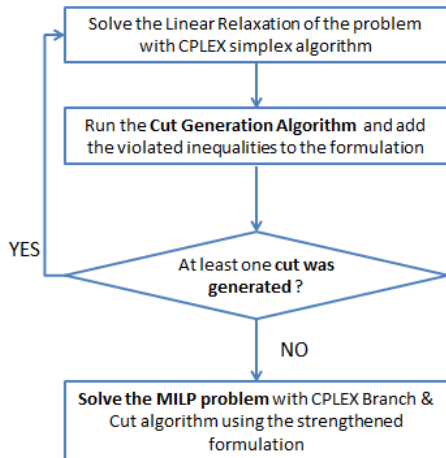
Separation

Polynomial in $O(N^2)$

Key point in practice

Careful selection of the valid inequalities to be added to the formulation

Cut & Branch algorithm



Preliminary computational results

Instances

Random generation based on the numerical values used by [Ahn *et al*, 2011]

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Numerical results

N	I	CPLEX alone			Cut & Branch algorithm			
		LP Gap	Opt Gap	Time	Cuts	LP Gap	Opt Gap	Time
105	5	4.7%	0.0%	2.4s	965	1.1%	0.0%	2.7s
200	5	7.2%	0.0%	9.6s	2241	1.2%	0.0%	7.5s
425	5	13.1%	0.0%	91s	4057	1.7%	0.0%	14s
605	5	14.6%	0.0%	522s	6438	1.5%	0.0%	33s
728	5	5.4%	0.0%	588s	4181	1.5%	0.0%	212s
1092	5	8.9%	0.0%	746s	8994	1.4%	0.0%	197s
1457	5	10.1%	0.1%	900s	14310	1.3%	0.0%	650s
1705	5	15.9%	0.2%	900s	17191	1.6%	0.1%	900s

- Average values for 5 randomly generated instances
- Resolution with CPLEX 12.6 on a PC running under Windows 10, Intel Core i7, 8 GB of RAM
- Time limit : 900s

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Conclusion and perspectives

Conclusion

- Lot sizing for a three-echelon remanufacturing system
- Uncertainty on returns and demand
- Multi-stage stochastic programming approach
- Polyhedral approach based on 'path inequalities'
- Significant reduction of the computation time

Conclusion and perspectives

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Perspectives

- Solving
 - More efficient cut generation for the 'path inequalities'
 - Extension to 'tree inequalities'
- Modeling
 - Heterogenous quality of the returns
 - Uncertainty on the disassembly yield

Thank you for your attention !